

Vellvm: Formalizing the Informal

Calvin Beck, Hanxi (Gary) Chen, Steve Zdancewic

University of Pennsylvania

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Thanks to these Collaborators!

Calvin Beck

Hanxi (Gary) Chen

Dmitri Garbuzov

Olek Gierczak

Paul He

Gil Hur

William Mansky

Milo M.K. Martin

Noé De Silva

Roger Burtonpatel

Santosh Nagarakatte

Gregory Malecha

Benjamin Pierce

Christine Rizkallah

Lucas Silver

Li-yao Xia

Irene Yoon

Yannick Zakowski

Jianzhou Zhao

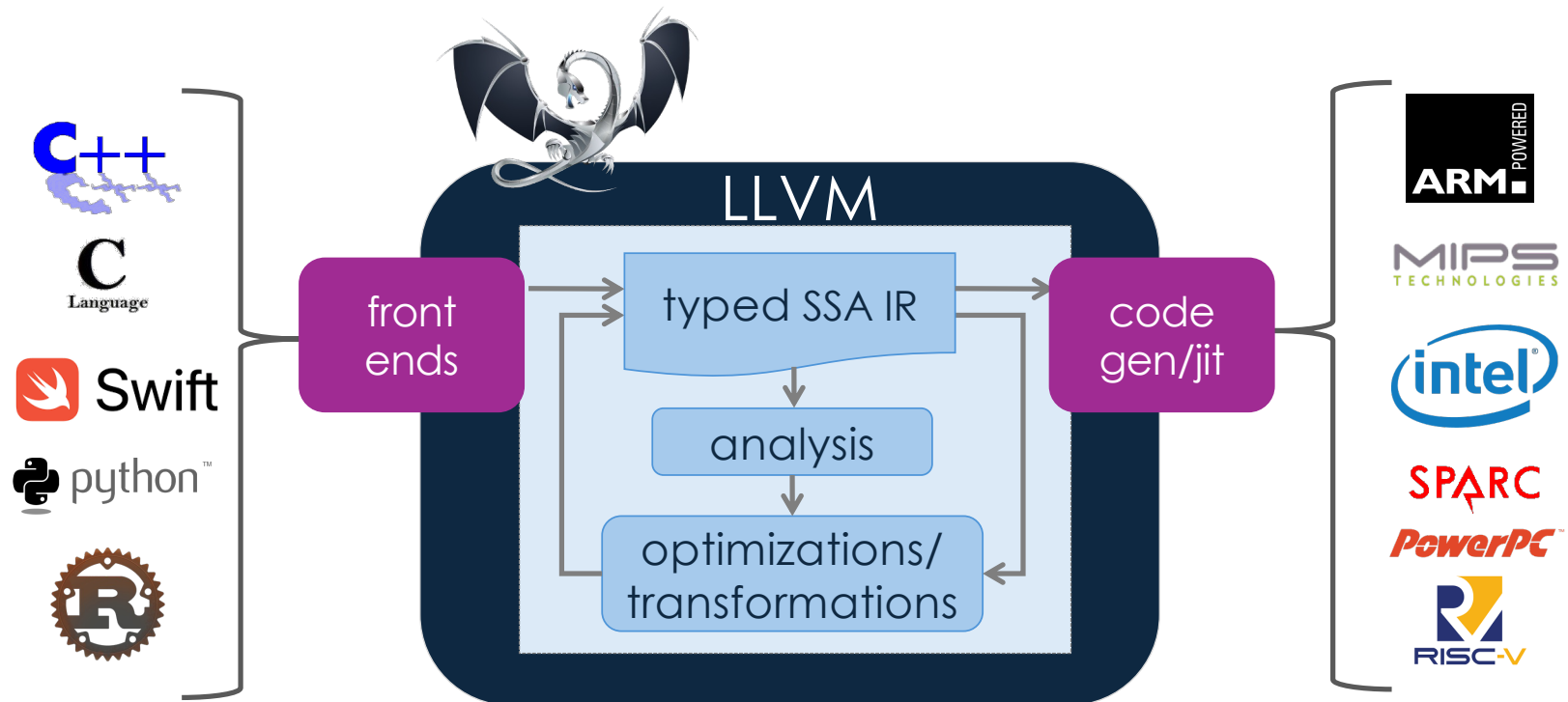


Outline

- **LLVM IR**
 - **pain points: undef ptrtoint/ lack of intptr type, varargs, platform-specific undef + memory**
- Vellvm
 - Sketch of Formalization
- Rocq Experience
- Differential Testing
 - GenLLVM structure
 - results
- Future

LLVM Compiler Infrastructure

[Lattner,2002]



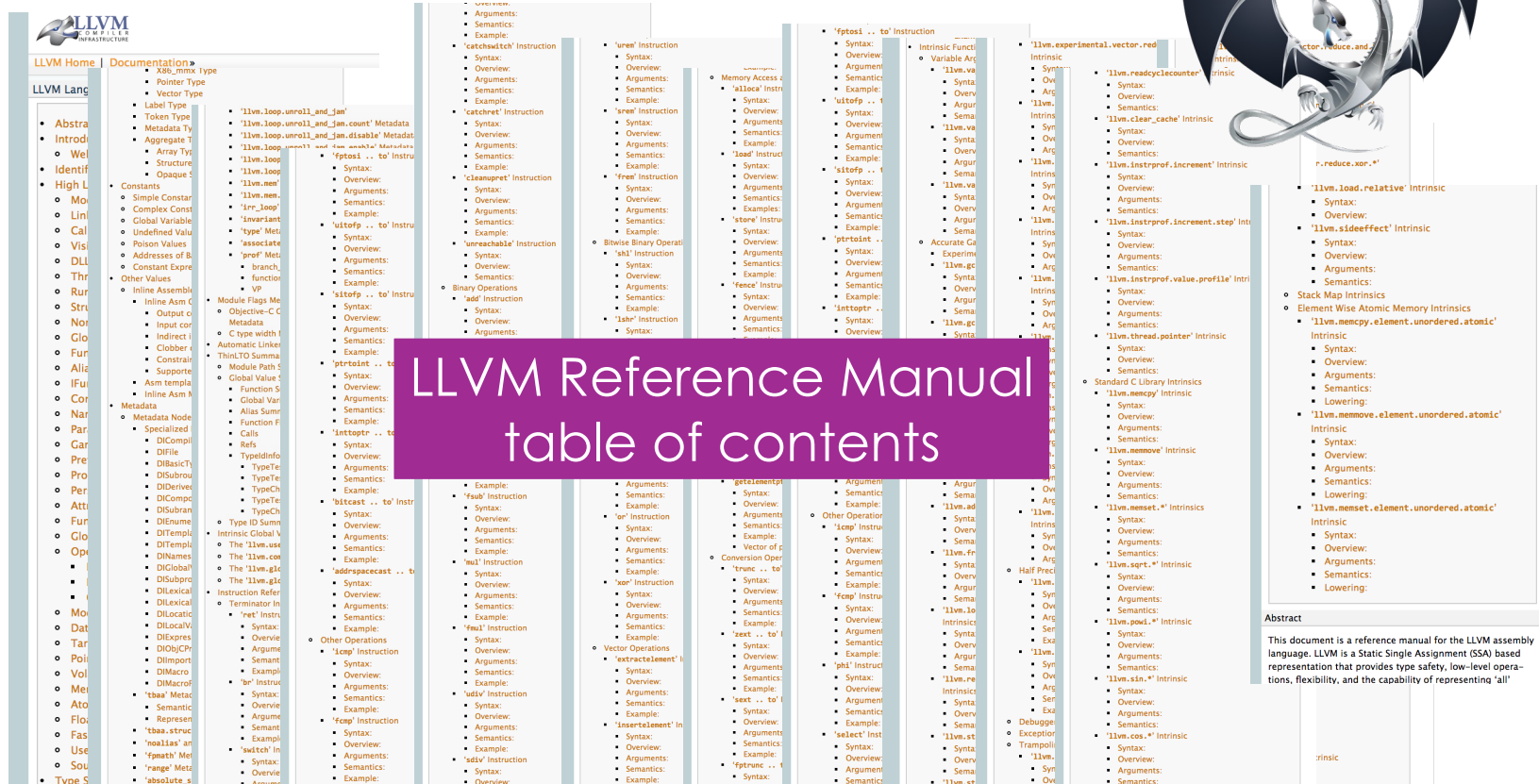
Translating C to LLVM IR

```
factorial(int n) {  
    int acc = 1;  
    while(n > 0) {  
        acc = acc * n;  
        n--;  
    }  
    return acc;  
}
```



```
define i64 @factorial(i64 %n){  
    %acc = alloca i64  
    store i64 1, i64* %acc  
    br label %start  
  
start:  
    %n1 = phi i64 [%n, %0], [%n2, %then]  
    %c = icmp sgt i64 %n1, 0  
    br i1 %c, label %then, label %end  
  
then:  
    %x1 = load i64, i64* %acc  
    %x2 = mul i64 %x1, %n1  
    store i64 %x2, i64* %acc  
    %n2 = sub i64 %n1, 1  
    br label %start  
  
end:  
    %ans = load i64, i64* %acc  
    ret i64 %ans  
}
```

But... it's complex



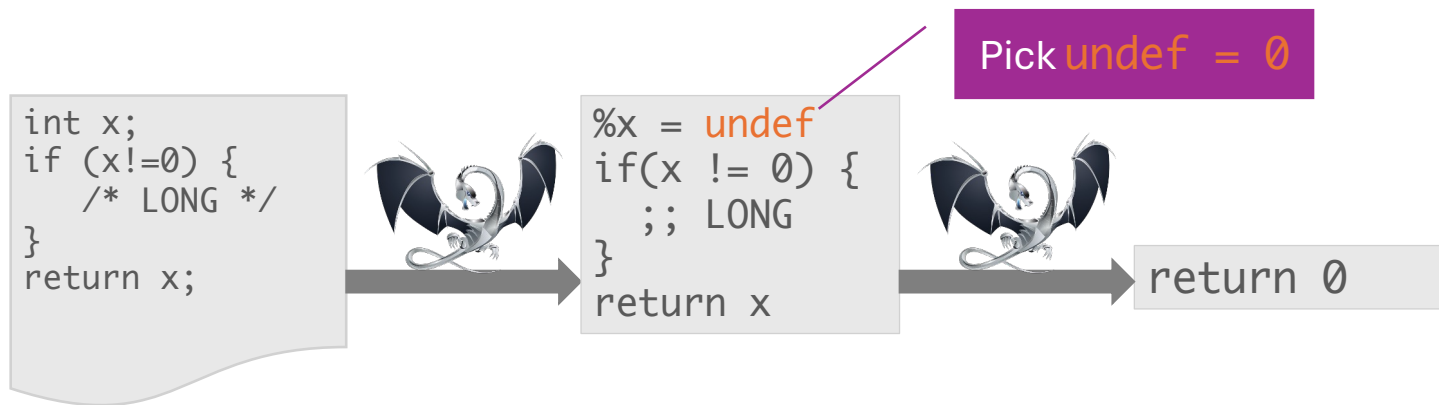
One Example: undef

The undef "value" represents an arbitrary, but indeterminate bit pattern for any type.

Used for:

- uninitialized registers
- reads from volatile memory
- results of some underspecified operations

Optimization with undef



(Slightly artificial example, branching on undef is UB in LLVM)

Optimistically use the best value
⇒ significantly faster code!

Interactions with Optimizations

Consider:

```
%y = mul i8 %x, 2
```

$[[\%x]] = [\text{i8 undef}]$
 $= \{0, 1, 2, 3, 4, 5, \dots, 255\}$
 $[[\%y]] = \{a \text{ mul } 2 \mid a \in [[\%x]]\}$
 $= \{0, 2, 4, \dots, 254\}$

versus:

```
%y = add i8 %x, %x
```

$[[\%x]] = [\text{i8 undef}]$
 $= \{0, 1, 2, 3, 4, 5, \dots, 255\}$
 $[[\%y]] = \{a + b \mid a \in [[\%x]], b \in [[\%x]]\}$
 $= \{0, 1, 2, 3, 4, \dots, 255\}$

≠

What are the Challenges?

Bug List: (12 of 435) [First](#) [Last](#) [Prev](#) [Next](#) [Show last search results](#)

Bug 33165 - Simplify* cannot distribute instructions for simplification due to undef

Status: REOPENED

Reported: 2017-05-25 02:13 PDT by Nuno Lopes

Davide Italiano 2017-05-25 08:55:40 PDT

[Comment 6](#)

Davide Italiano 2017-05-25 09:05:26 PDT

[Comment 7](#)

(unless we want to give up on some undef transformations, and special case select, but I'm afraid others might be affected too)

John Regehr 2017-05-25 09:09:24 PDT

[Comment 8](#)

Yes, this is one of those test cases. There are so many optimization failures that Nuno has been automatically filtering out classes of mistranslation that are known to be hard to fix but I guess he decided to take a closer look at some of them.

Soon I'll be able to include branches/phis in these test cases, but only forward branches due to a limitation in Alive.

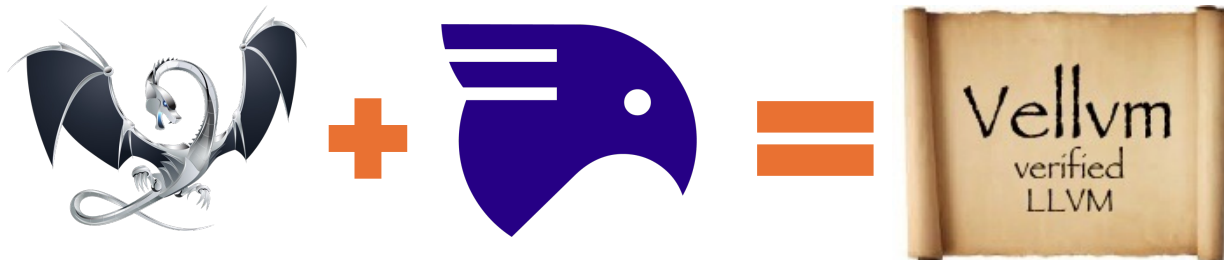
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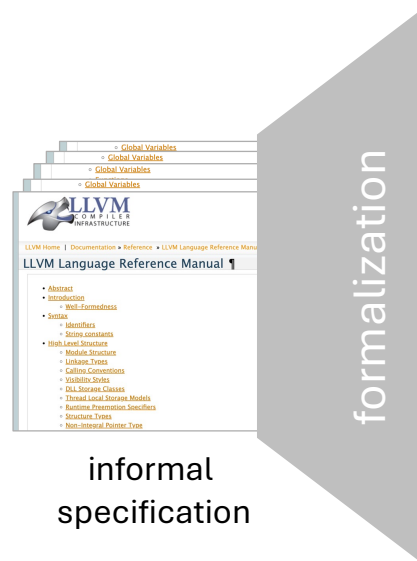
What Can We Do?

- LLVM is important
- But complex, features interact subtly
- Can lead to bugs :(

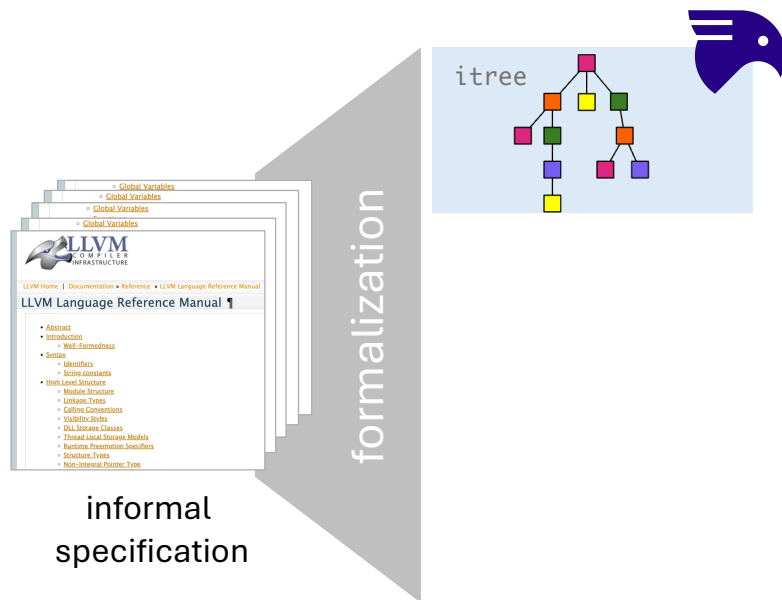
Maybe we can formalize it?



Structure of the Formalism: Informal Spec

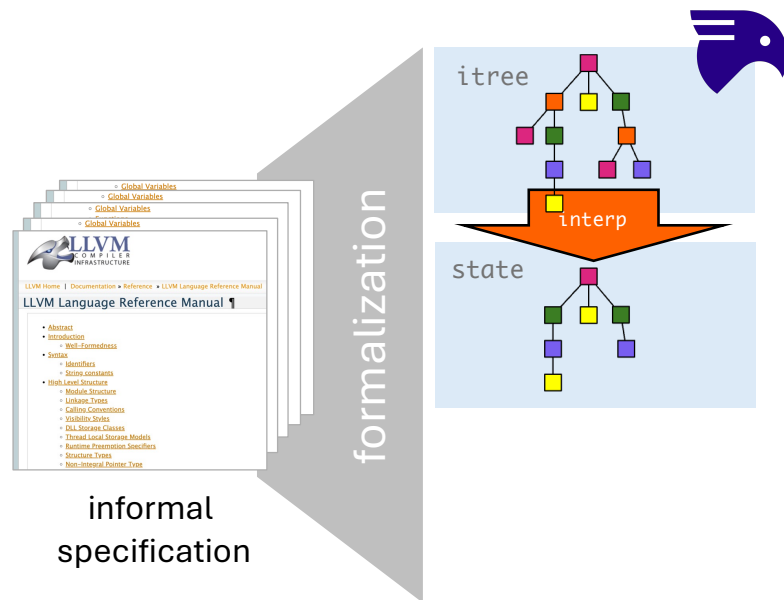


Structure of the Formalism: Control Flow



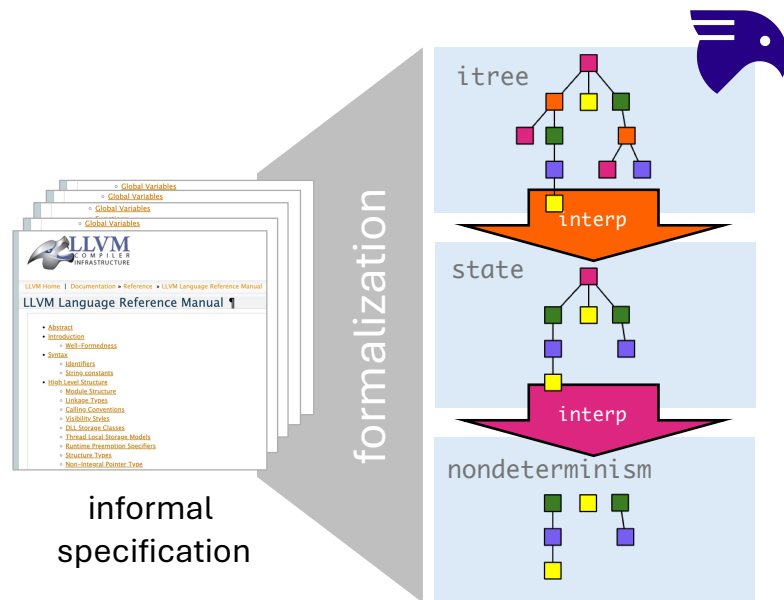
- Abstract Syntax
- Interaction Trees
 - "free monad" (nearly syntax)
 - infinitary behaviors

Structure of the Formalism: Monadic Interpreters

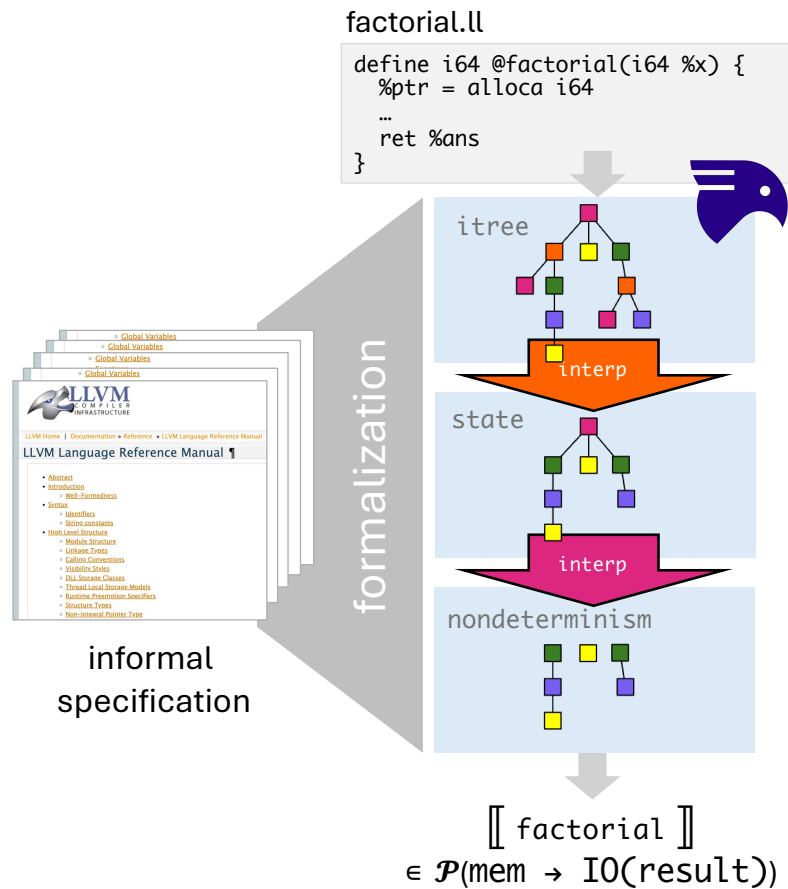


- Abstract Syntax
- Interaction Trees
 - "free monad" (nearly syntax)
 - infinitary behaviors
- Modular Event Handlers
 - Interpret the events
 - e.g., to introduce state

Structure of the Formalism: Nondeterminism



- **Abstract Syntax**
- **Interaction Trees**
 - "free monad" (nearly syntax)
 - infinitary behaviors
- **Modular Event Handlers**
 - Interpret the events
 - e.g., to introduce state
- **Nondeterminism**
 - Interpret into Rocq's Prop



Net Result

1. End-to-end specification

- defines sets of **allowed** behaviors

2. Prove refinements *between* layers

- equivalence earlier implies equivalence later

Status of Vellvm

- Large subset of *sequential* LLVM IR
 - Most C programs, modulo libraries
 - A few "intrinsic" + support for adding more
- Rich memory model [ICFP24]
- Verified interpreter with some nice features
 - `printf`
 - Command line argument support for LLVM programs
- GenLLVM $\Rightarrow$ randomized testing

Upshot: can interpret many compiled C programs, depending on library usage.

Relatively easy to add new intrinsics and library functions.

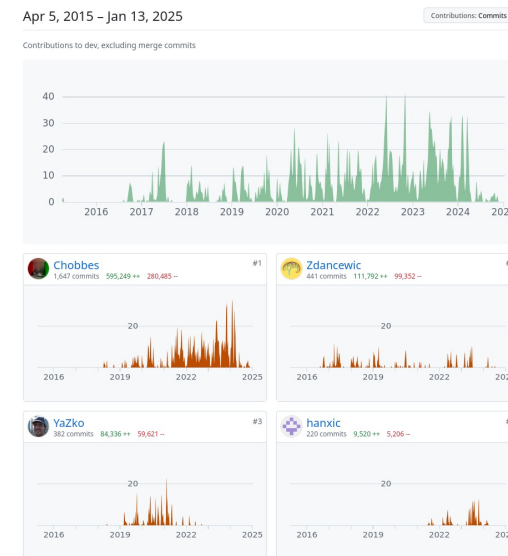
(Rough) Statistics about Vellvm Development

Lines of Spec	Lines of Proof
51.4K	95.6K

Size of vo files	Gzipped vo file size
516M	259M

Libraries:

- flocc
- paco
- itrees
- extlib
- quickchick
- (menhir)



Benefits of formalizing LLVM IR

Finding Bugs/Ambiguities

- Size 0 type issues
- sub-byte-sized pointer indices
- address guessing
 - "Reconciling high-level optimizations and low-level code in LLVM" [OOPSLA 18]
- lack of intptr_t (see ICFP 24)
- undef "entanglement" (see ICFP 24)
- ...

Ambiguities in LLVM IR

"Allocating zero bytes is legal, but the **returned pointer may not be unique**" -- LangRef

"The allocated memory is uninitialized, and **loading from uninitialized memory produces an undefined value**" -- LangRef

```
%p = alloca [0 x i32]  
%x = load [0 x i32], [0 x i32]* %p
```

Is %x undef?

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Experience

Pros:

- Consistent rich logic
- Great for finding bugs / ambiguities
- Modules can be great
- Extraction (works shockingly well)

Cons:

- Constraints (Coinductive + existentials)
- Module issues
- More first class module types?
- Extraction bugs
- Standard Library
- Compilation time
- Labor intensive



Rocq Experience: Coinductive Existentials

- Arises due to nondeterminism
- Many constraints.

Rocq Experience: Coinductive Existentials

- Arises due to nondeterminism
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```
a : itree A X
b : itree B X
b2 : itree B2 X
R : a ≈ b (* Coinductive *)
MODEL : b2 ∈ modelB b
-----
∃ (a2 : itree A2 X),
  a2 ∈ modelA a ∧
  a2 ≈ b2.
```



<https://github.com/Chobbes/existential-coinduction-experiments>

Outline

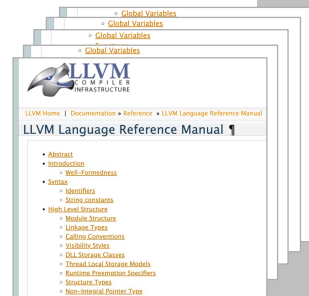
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Validity of Vellvm

- We have a large **formal semantics** for LLVM IR
 - Formal interpreter using interaction trees and a two-phase infinite/finite memory model.
- But **does it adhere to the LLVM Language Reference?** 🤔
 - We want Vellvm semantics faithfully model the informal specification.
- Testing!
 - Unit Testing (manual + Alive2)
 - **Differential Testing** (like CSmith and YARPGen)

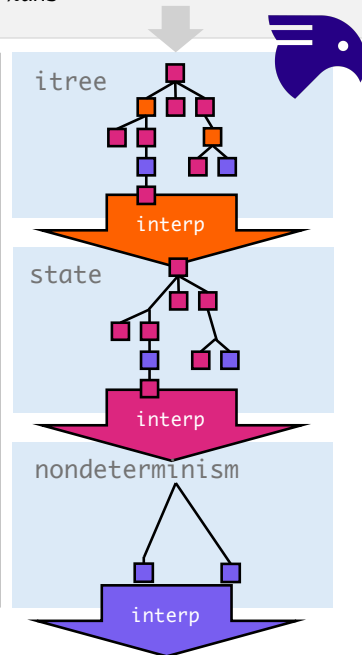
factorial.ll

```
define i64 @factorial(i64 %x) {
  %ptr = alloca i64
  ...
  ret %ans
}
```

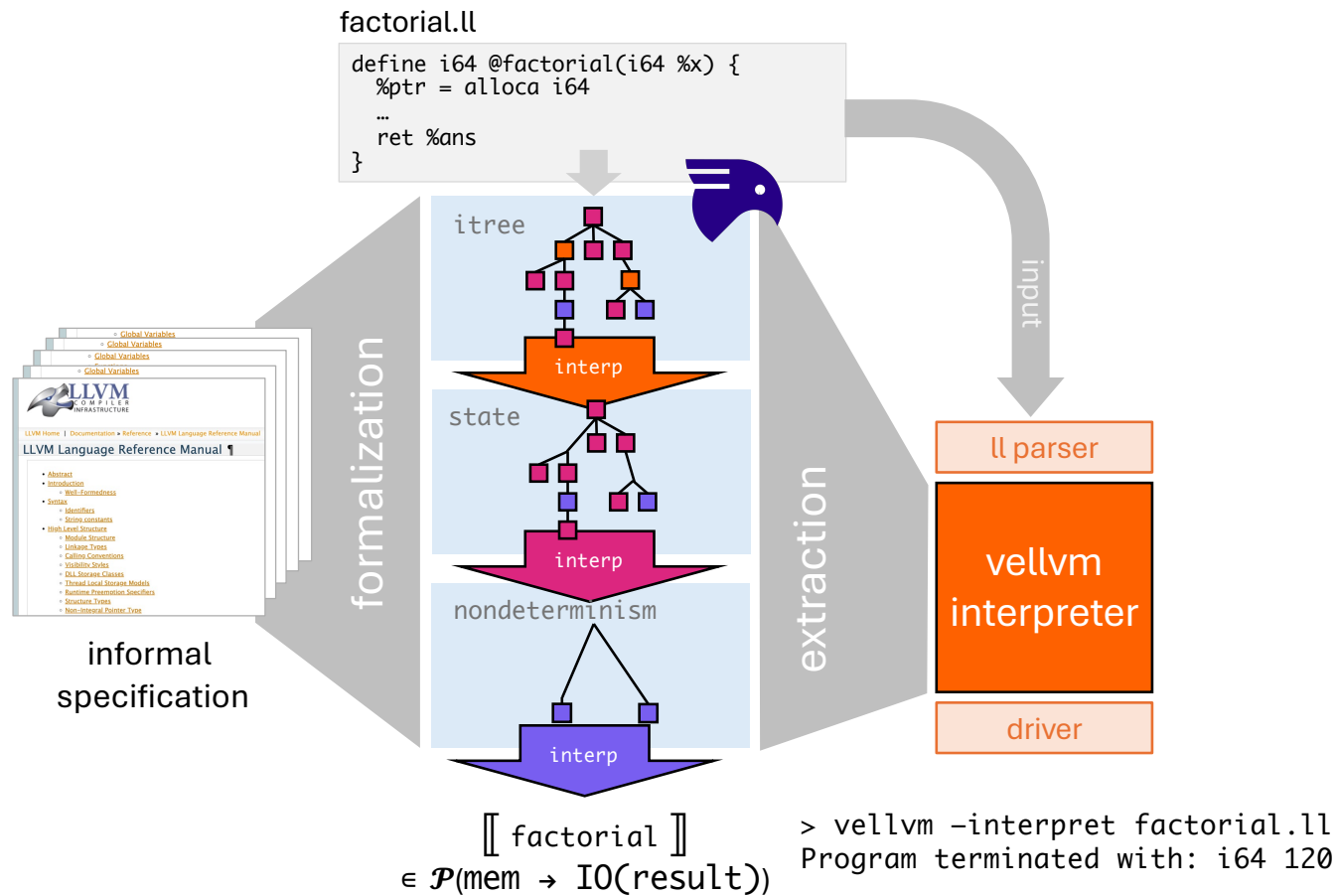


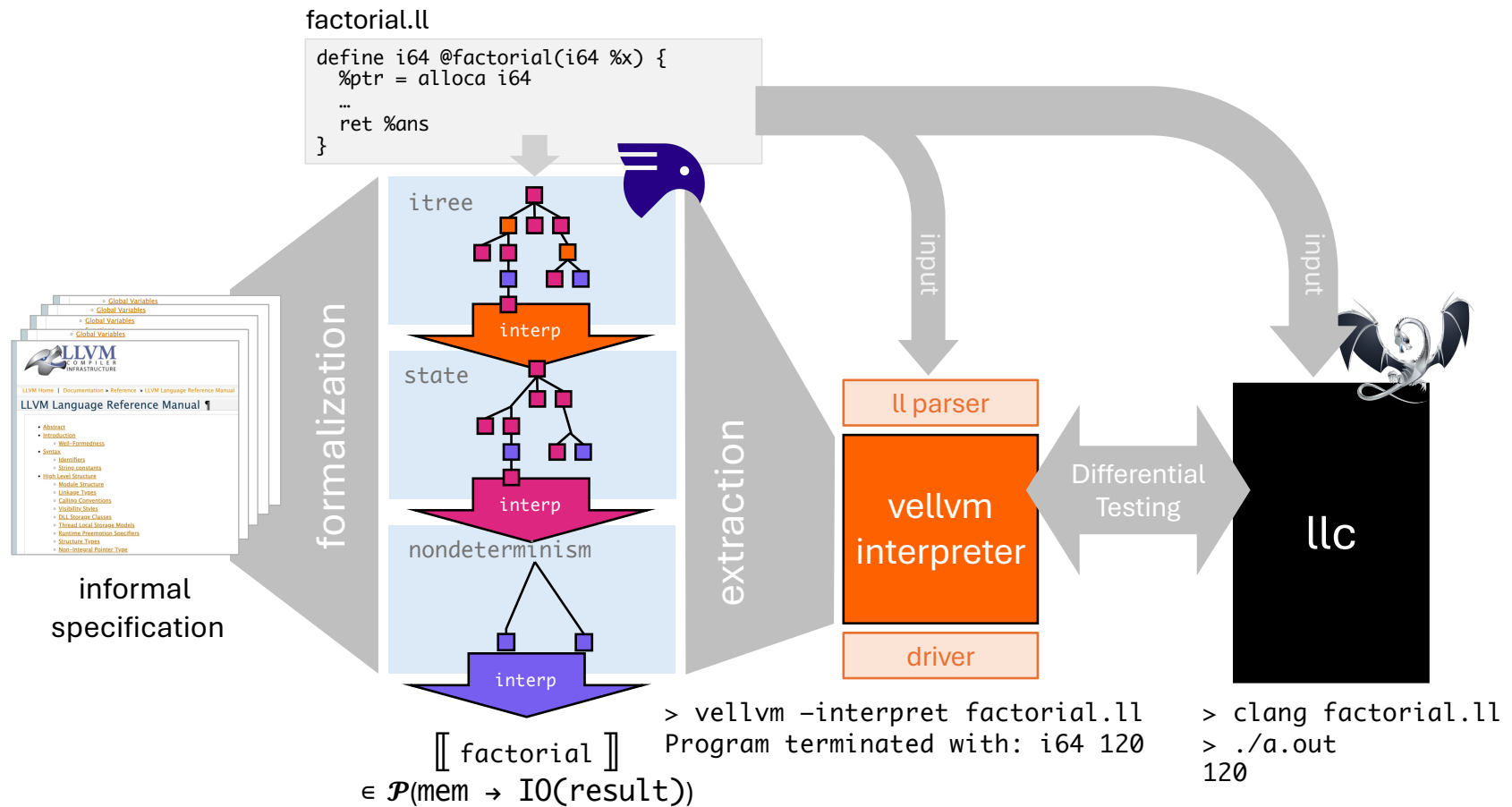
informal
specification

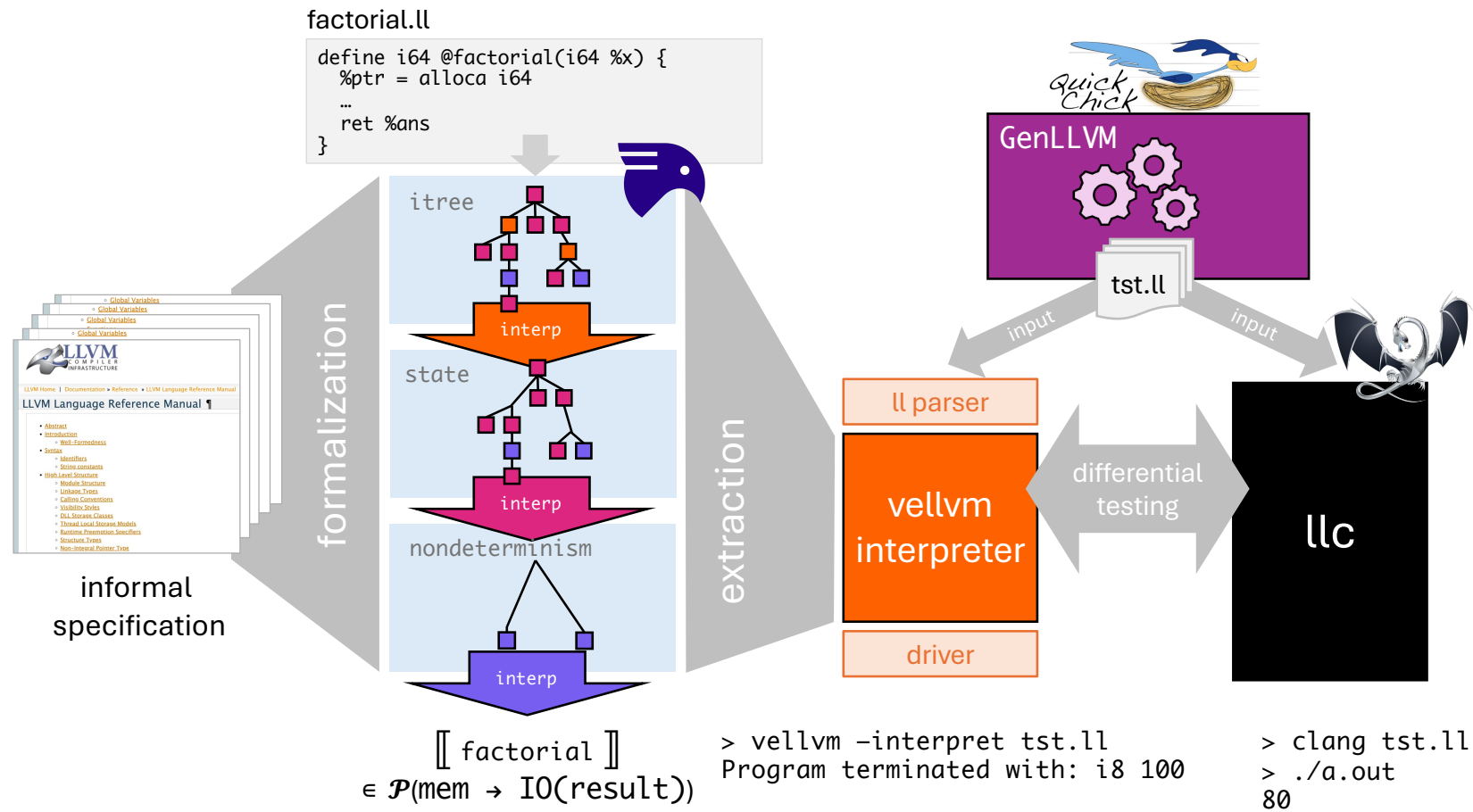
formalization



$\llbracket \text{factorial} \rrbracket$
 $\in \mathcal{P}(\text{mem} \rightarrow \text{IO}(\text{result}))$







Trickiness in Differential Testing of LLVM IR

Requirements for LLVM Programs

- Deterministic
- Well-formed
 - Type-safe, well-scoped, etc.
- Undefined-behavior free

Desirable Properties for Testing

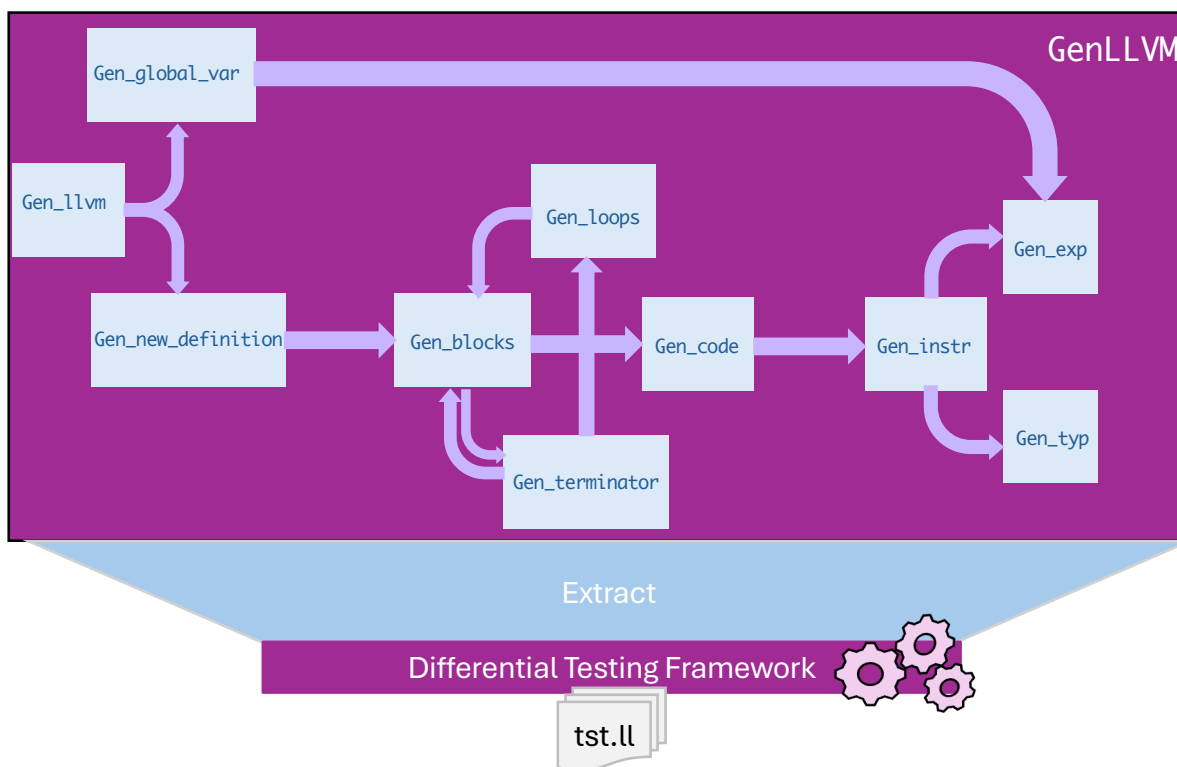
- Exercise large subset of LLVM syntax
- Nontrivial control flows
- Random code patterns

GenLLVM: Infrastructure



- Built using the G monad from the QuickChick library
- The GenLLVM monad can
 - Early termination when generating ill-formed code
 - Use a context-aware metadata system (inspired by Ecstasy)
 - Generate types / expressions / terminators

GenLLVM: Generation Strategy



Algorithm 1

```

procedure GEN_TERMINATOR( $t, \text{back\_blocks}, \text{SIZE}$ )  $\triangleright$  Generate an LLVM Function
     $r \leftarrow \text{RANDOM}(0, \epsilon_1 + \epsilon_2 + \epsilon_3 + \epsilon_4 + \epsilon_5)$   $\triangleright$  Each  $\epsilon_i$  is a tunable frequency
    if  $r \leq \epsilon_1$  then  $\triangleright$  Generate return
        return ( $\text{gen\_ret } t$ )
    else if  $r \leq \epsilon_2$  then  $\triangleright$  Simple branching
        ( $b, \text{blocks}$ )  $\leftarrow \text{gen\_blocks}(t, \text{back\_blocks}, \frac{\text{SIZE}}{2})$ 
        return ( $b, t, b, \text{blocks}$ )
    else if  $r \leq \epsilon_3$  then  $\triangleright$  Conditional branching
        ( $b_1, \text{blocks}_1$ )  $\leftarrow \text{backtrack}(\text{gen\_blocks}(t, \text{back\_blocks}, \frac{\text{SIZE}}{2}))$ 
        ( $b_2, \text{blocks}_2$ )  $\leftarrow \text{gen\_blocks}(t, \text{back\_blocks}, \frac{\text{SIZE}}{2})$ 
         $c \leftarrow \text{gen\_exp\_size}(0, 1)$ 
        return ( $b_1, 1, c, b_1, b_2, \text{blocks}_1 + \text{blocks}_2$ )
    else if  $r \leq \epsilon_4$  then  $\triangleright$  Generate Loop
        ( $t, b, \text{blocks}$ )  $\leftarrow \text{gen\_loop}(t, \text{back\_blocks}, \frac{\text{SIZE}}{2})$ 
         $b' \leftarrow \text{add } b \text{ to the head of } \text{blocks}$ 
        return ( $t, b'$ )
    else if  $r \leq \epsilon_5$  then  $\triangleright$  Generate recursion
         $b \leftarrow \text{one\_of}(\text{back\_blocks})$ 
        return ( $b, t, b, []$ )

```

Figure: Algorithm for generating terminator

```

gen_typ : GenLLVM typ
gen_exp : typ  $\rightarrow$  GenLLVM exp
gen_instr : GenLLVM (list instr_id  $\times$  instr)
gen_code : GenLLVM (list instr_id  $\times$  instr)
gen_ret : typ  $\rightarrow$  GenLLVM terminator
gen_terminator : typ  $\rightarrow$  list block_id  $\rightarrow$ 
    GenLLVM (terminator  $\times$  list block)
gen_blocks : typ  $\rightarrow$  list block_id  $\rightarrow$ 
    GenLLVM (block  $\times$  list block)
gen_loop : typ  $\rightarrow$  list block_id  $\rightarrow$ 
    N  $\rightarrow$  GenLLVM (block  $\times$  list block)
gen_new_definition : typ  $\rightarrow$  list typ  $\rightarrow$  GenLLVM definition
gen_global_var : GenLLVM global
gen_llvm : GenLLVM (list toplevel_entity)

```

Figure: Types of GenLLVM's functions

Some Trickiness: inttoptr

```
%x = add i32 123, 0  
%y = inttoptr i32 %x to i8*  
%z = load i8, i8* %y  
ret %z
```

- **Nondeterminism**
- Csmith and YARPGen do not generate `inttoptr` (mostly)!
- **Observation:** `inttoptr` can be invoked on integers that from `ptrtoint`!
 - (following size restrictions)

Evaluation

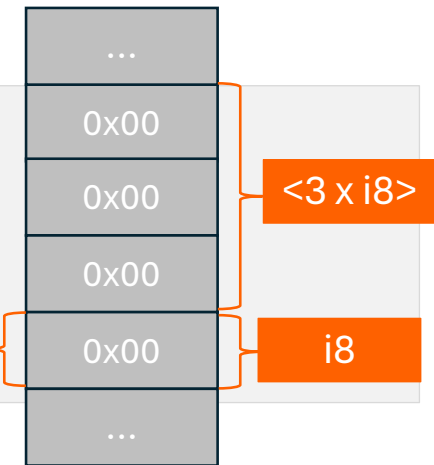
✓✓: emphasized; ✓: generated;
~: constrained; × : not generated

Feature	Csmith	YARPGen V1	YARPGen V2	GenLLVM
Level	C	C / C++	C / C++	LLVM
Int Arith	✓✓ (wrappers)	✓	✓	✓
Float Arith	✓	~ (constrained)	~ (constrained)	~ (LLVM Semantics)
Loops	✓ (induction var.)	✓	✓✓	✓ (induction var.)
Vectors	×	~ (short vectors)	✓✓	✓
Agg. Data	✓ (C)	✓ (C)	✓ (C)	✓✓
Ptr. Casts	✓ (only to void*, C)	×	×	✓

- Ran Vellvm on 10,000 programs each from GenLLVM, Csmith, YARPGen versions 1 and 2
 - (extracted) Vellvm can execute 2.5M+ LLVM IR instructions in 2m03s.
 - ✓: generated; ~: constrained; × : not generated

We Found Bugs in Clang!

```
define i8 @main() {  
  %v0 = alloca i32  
  store i32 0, i32* %v0, align 1  
  %v1 = ptrtoint i32* %v0 to i64  
  %v2 = inttoptr i64 %v1 to <{<3 x i8>, i8}>*  
  %v3 = load <{<3 x i8>, i8}>, <{<3 x i8>, i8}>* %v2, align 1  
  %v4 = extractvalue <{<3 x i8>, i8}> %v3, 1  
  ret i8 %v4  
}
```



- "zero-initialize" an **i32** into a **32-bit packed-struct**
- Read the **last byte** from it
- Semantically, 0 should be returned!

But: clang: 46
 Vellvm: 0

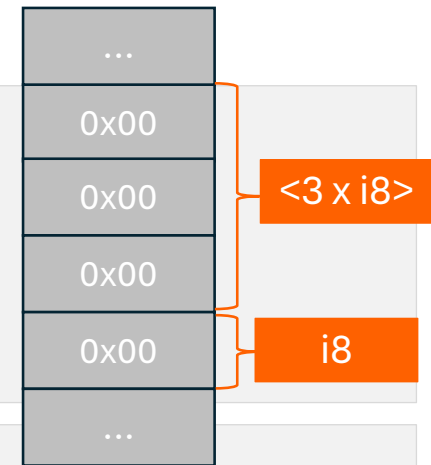
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    %v2 = inttoptr i64 %v1 to <{T, i8}>*  
    %v3 = load <{T, i8}>, <{T, i8}>* %v2, align 1  
    %v4 = extractvalue <{T, i8}> %v3, 1  
    ret i8 %v4  
}
```

clang -S

T = <3 x i8>

```
main:  
    addi    sp, sp, -16  
    mv      a0, zero  
    sw      a0, 12(sp)  
    lb      a0, 16(sp)  
    addi    sp, sp, 16  
    ret
```



- Off-by-one error in extractvalue with packed struct with vector

We Found Bugs in Clang!

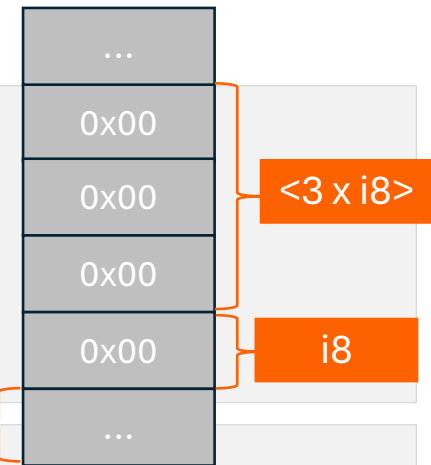
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```

clang -S

return



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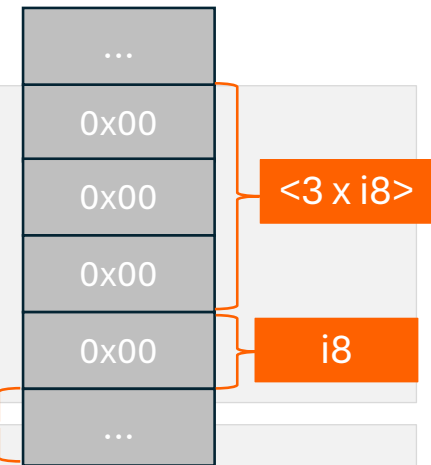
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```

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    addi    sp, sp, -16  
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```

clang -S

return



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    %v1 = ptrtoint i32* %v0 to i64  
    %v2 = inttoptr i64 %v1 to <{T, i8}>*  
    %v3 = load <{T, i8}>, <{T, i8}>* %v2, align 1  
    %v4 = extractvalue <{T, i8}> %v3, 1  
    ret i8 %v4  
}
```

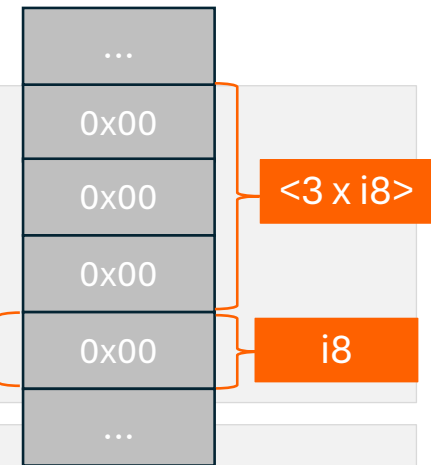
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    lb      a0, 16(sp)  
    addi    sp, sp, 16  
    ret
```

clang -S

T = [3 x i8]

```
main:  
    addi    sp, sp, -16  
    mv      a0, zero  
    sw      a0, 12(sp)  
    lb      a0, 15(sp)  
    addi    sp, sp, 16  
    ret
```



- Off-by-one error in extractvalue with packed struct with vector

Takeaways

- Rocq is great for formalizing complex systems!
- Coinduction & ITrees work well for defining semantics
- Executability & differential testing were essential
 - for debugging our semantics
 - for finding bugs in Clang & Langref

Together: these make Vellvm a useful platform for understanding LLVM IR semantics

Wrap-up / Ongoing / Future Work

- Better I/O / C library support
- Factoring out the memory model
- Verified Analyses and Optimizations
- Concurrency
 - "A concurrency model based on monadic interpreters"
[Chappe et al. CPP25]
- LLVM IR version bump
- More instructions in GenLLVM



<https://github.com/Vellvm/vellvm>



Appendix

Rocq Experience: Extraction

- Shockingly Robust
- Nice features like Extraction Blacklist, etc.
- A little buggy at times
 - Modules?
 - Things named 'int'?

Rocq Experience: Modules

- Modules / module types are fantastic!
- Wish they were a bit more first-class
 - Construct anonymous module types with ($\leq +$) operator?
- Careful with Program in Module Types...

Undef Interactions with Memory

Important for optimizations like *store forwarding*

```
%x = i16 select i1 undef, i16 0x1234, i16 0x0000  
%y = i16 select i1 undef, i16 0x5678, i16 0x0000  
store i16 %x, ptr %ptr  
store i16 %y, ptr (inttoptr (add iptr (ptrtoint %ptr) 2)  
%z = load i32, ptr %ptr
```

Comparison of Instructions Exercised

